### Introduction to big Ramsey degrees

Part 1: Big Ramsey degrees of rationals and Rado graph

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I would like to cover some old&new results in the area of big Ramsey degrees:

- 1 Big Ramsey degrees of rationals and Rado graph.
- 2 Recent progress in the area
  - 1 Big Ramsey degrees of Triangle-free graphs
  - 2 New Ramsey theorem for trees with successor operation.
- 3 Applications of the new Ramsey theorem
  - 1 Easy proof of unrestricted Nešetřil-Rödl or Abramson-Harington theorem
  - **2** Big Ramsey degrees of structures forbidding bigger substructures.

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$$\forall_{p,k\geq 1}:\omega\longrightarrow (\omega)_{k,1}^p.$$

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Definition (Erdős–Rado partition arrow)

 $N \longrightarrow (n)_{k,t}^{p}$  means: For every partition of  $\binom{N}{p}$  into *k* classes (colours) there exists  $X \in \binom{N}{n}$  such that  $\binom{X}{p}$  belongs to at most *t* parts.

 $(t = 1 \text{ means that } \begin{pmatrix} \chi \\ \rho \end{pmatrix}$  is monochromatic.)

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In 1970's a concept of structural Ramsey theory was introduced. A Ramsey theorem can be seen as a theorem about the class of linear orders.

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Let  $\mathcal{O}$  be the class of all finite linear orders.

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 $\begin{pmatrix} B \\ A \end{pmatrix}$  is the set of all embeddings of structure **A** to structure **B**.

Definition (Leeb's generalization of the Erdős-Rado partition arrow)

 $\mathbf{C} \longrightarrow (\mathbf{B})_{k,t}^{\mathbf{A}}$  means:

For every *k*-colouring of  $\binom{\mathbf{C}}{\mathbf{A}}$  there exists  $f \in \binom{\mathbf{C}}{\mathbf{B}}$  such that  $\binom{f(\mathbf{B})}{\mathbf{A}}$  has at most *t* colours.

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A natural question: Is the same true for  $(\mathbb{Q}, \leq)$  (the order of rationals)?

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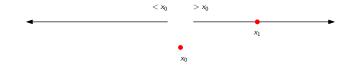
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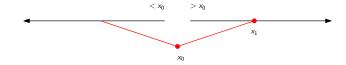
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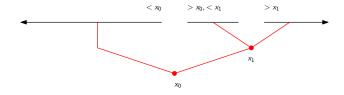
Sierpiński: not true for |O| = 2.

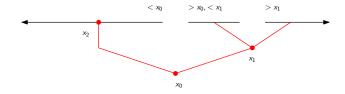


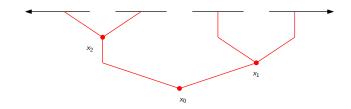


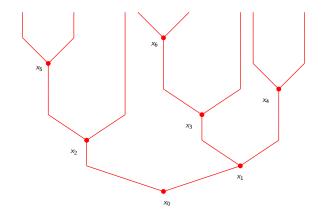


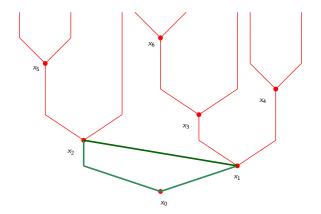


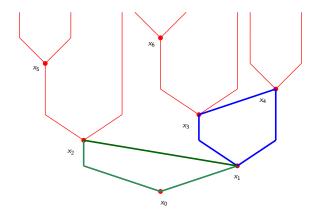


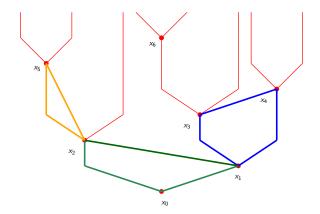


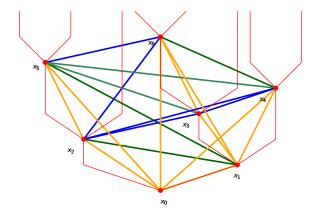












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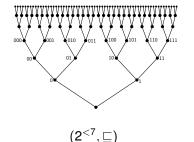
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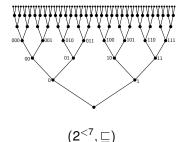
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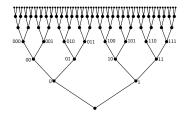
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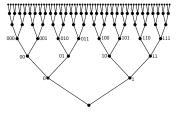


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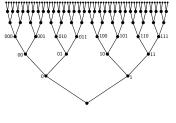
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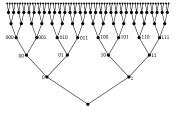
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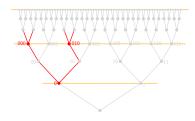
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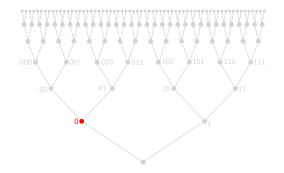


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- The height of *T*, denoted by h(T), is the minimal natural number *h* such that  $T(h) = \emptyset$ . If there is no such number *h*, then we say that the height of *T* is  $\omega$ .

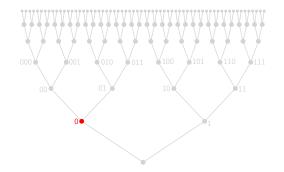


- A subtree of a tree *T* is a subset *S* ⊆ *T* viewed as a tree equipped with the induced partial ordering.
- Given a tree *T* and nodes *s*, *t* ∈ *T* we say that *s* is a successor of *t* in *T* if *t* ≤<sub>*T*</sub> *s*.
- The node s is an immediate successor of t in T if t <<sub>T</sub> s and there is no s' ∈ T such that t <<sub>T</sub> s' <<sub>T</sub> s.
- Node with no successors is leaf.



Let *T* be rooted tree. Nonempty  $S \subseteq T$  is a strong subtree of *T* of height  $n \in \omega + 1$  if:

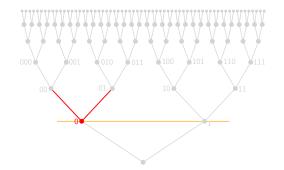
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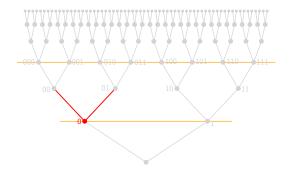


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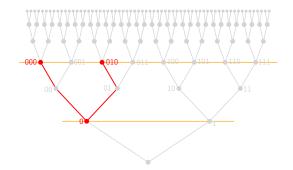


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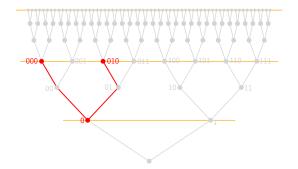


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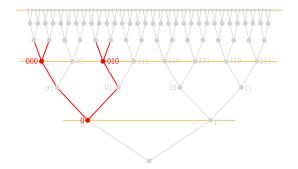


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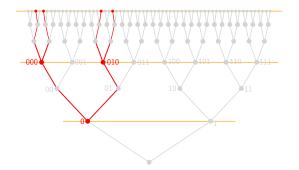


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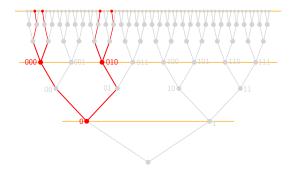


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- 4 S has height n.

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Theorem (Milliken 1979)

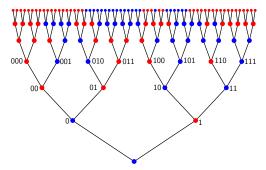
For every rooted finitely branching tree T with no leaves, every  $k \in \omega$  and every finite colouring of  $\operatorname{Str}_k(T)$  there is  $S \in \operatorname{Str}_\omega(T)$  such that the set  $\operatorname{Str}_k(S)$  is monochromatic.

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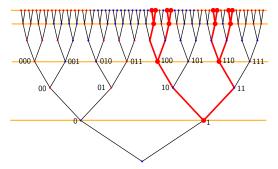


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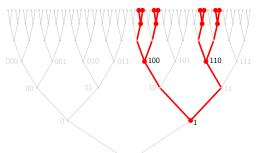


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Notice that for regularly branching tree the strong subtree is isomorphic to the original tree.

We aim to prove:

Theorem (Laver, late 1969)

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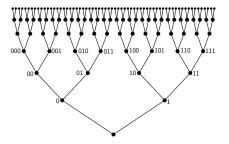
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Easy case |O| = 1:

 Nodes of the binary tree 2<sup><ω</sup> ordered "from left to right" yields (Q, ≤).

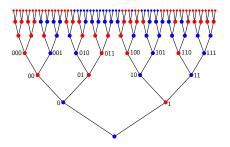


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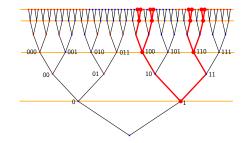


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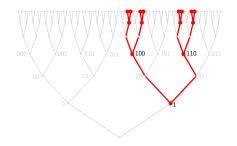


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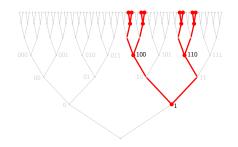
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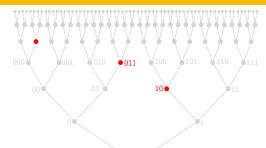
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Easy case |O| = 1:

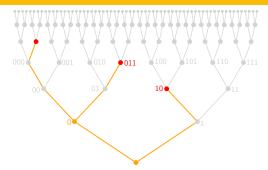
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If |O| = n > 1 we transfer colourings of *n*-tuples of nodes to colouring of strong subtrees.

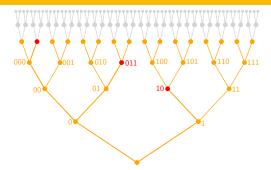




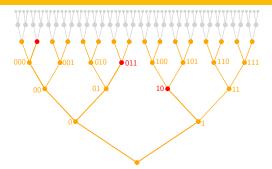
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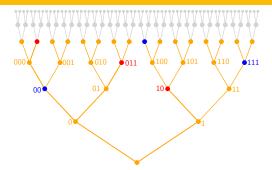
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Multiple choices of X may lead to a same envelope. We speak of different embedding types within a given envelope.

Now we can finish proof of:

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• Fix  $(O, \leq_O) \in \mathcal{O}$  and put n = |O|.

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- **5** The resulting copy will have at most T(n) different colours

# Big Ramsey degrees of $(\mathbb{Q}, \leq)$ are finite!



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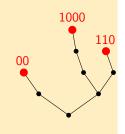
● Leaf: There exists precisely one  $a \in A$  with  $|a| = \ell$ . Moreover for every  $b \in A$ ,  $|b| > \ell$  it holds that  $b(\ell) = 0$ .

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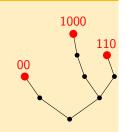
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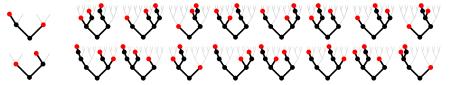
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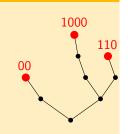




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Fun fact: Number of Devlin types of size n is

$$t_n = \sum_{\ell=1}^{n-1} {\binom{2n-2}{2\ell-1}} t_{\ell} \cdot t_{n-\ell}$$
 with  $n_1 = 1$ 

This is well known sequence (of the odd tangent numbers).

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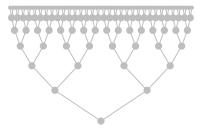
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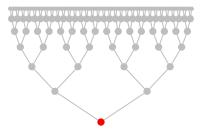
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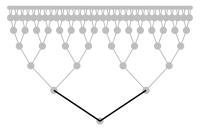
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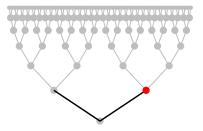
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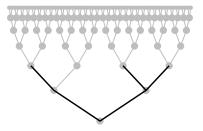
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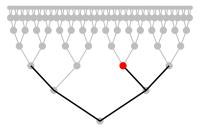
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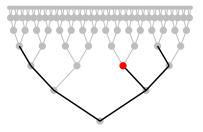
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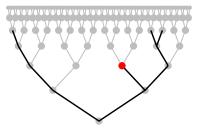
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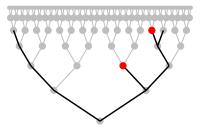
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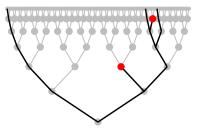
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Recall that in Devlin type on every is either leaf or branching. Minimal envelope will include all those levels where branching or leaf of *B* happens and skip all others.

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We thus obtain an upper bound: T(|O|) is at most the number of Devlin types.

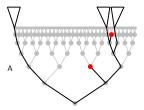
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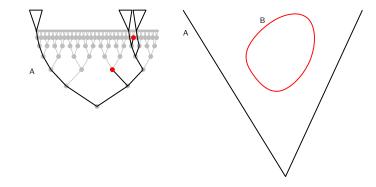


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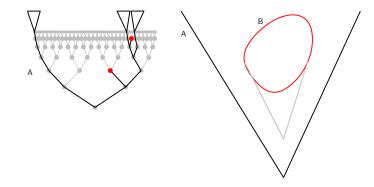


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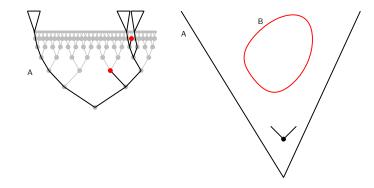


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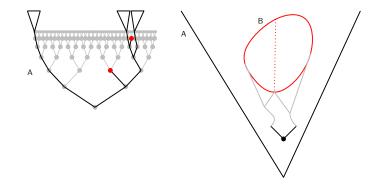


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Proceed by induction on the individual branching/leaf events of *A*. The meet splits leafs of *B* into two infinite intervals. Each with meet above its son.

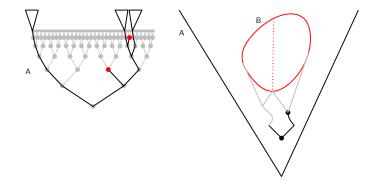


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Proceed by induction on the individual branching/leaf events of *A*. Use corresponding meet to realize 2nd branching of *A*.

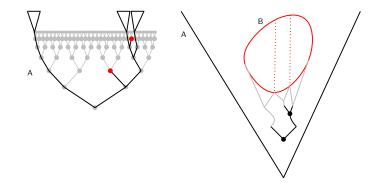


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Proceed by induction on the individual branching/leaf events of *A*. *B* further subdivides into intervals.

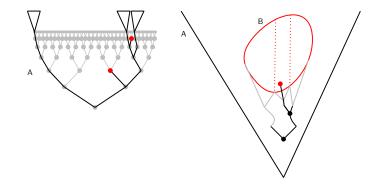


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### Proof by a slide-show.

Proceed by induction on the individual branching/leaf events of *A*. Continue analogously with next branching.

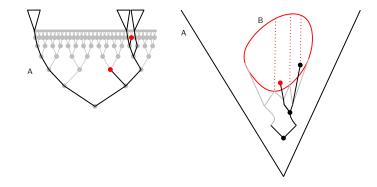


#### Lemma

Let A be a Devlin type representing  $(\mathbb{Q}, \leq)$  and  $B \subseteq A$  a copy of  $(\mathbb{Q}, \leq)$  then there exists a  $C \subseteq B$  whose embedding type (inside a minimal envelope) is A.

### Proof by a slide-show.

Proceed by induction on the individual branching/leaf events of *A*. Place the first leaf of *A* into corresponding interval.



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## Proof by a slide-show.

Proceed by induction on the individual branching/leaf events of *A*. ... blablabla... proof done.



We characterised the big Ramsey degrees of rationals and gave a closed-form formula. This shows that the upper bound proof is best possible. It also has applications to topological dynamics

Kechris–Pestov–Todorčević: Fraïssé limits, Ramsey theory, and topological dynamics of automorphism groups Zucker: Big Ramsey degrees and topological dynamics

Definition

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- We denote by  $\mathcal{G}$  the class of all finite graphs.

#### Theorem

$$\forall_{\mathbf{A}\in\mathcal{G}}\exists_{\mathcal{T}=\mathcal{T}'(\mathbf{A})\in\omega}\forall_{k\geq 1}:\mathbf{R}\longrightarrow(\mathbf{R})_{k,\mathcal{T}}^{\mathbf{A}}.$$

This theorem was published by Sauer in 2006 and also appears in Todorčević' Introduction to Ramsey spaces. Values of  $T'(\mathbf{G})$  were characterised by Laflamme–Sauer–Vuksanović in 2010.

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A finitary version is (probably more) famous!

Theorem (Nešetřil-Rödl 1977, Abramson-Harington 1978)

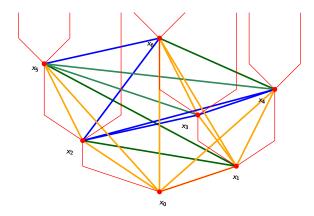
$$\forall_{\mathbf{A}\in\mathcal{G}}\exists_{t=t(\mathbf{A})\in\omega}\forall_{\mathbf{B}\in\mathcal{G},k\geq 1}\exists_{\mathbf{C}}\in\mathcal{G}:\mathbf{C}\longrightarrow(\mathbf{B})_{k,t}^{\mathbf{A}}.$$

## Understanding the unavoidable colourings

While trying to formulate Ramsey-type theorem it is good to check if there are any unavoidable colourings and if so understand their structure.

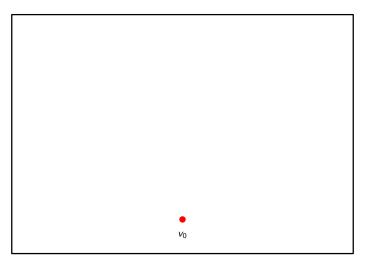
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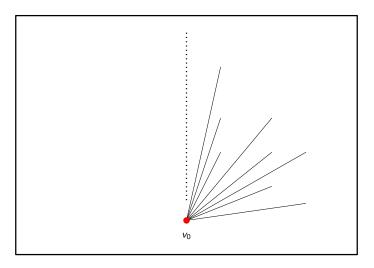
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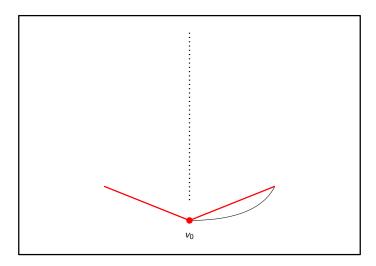


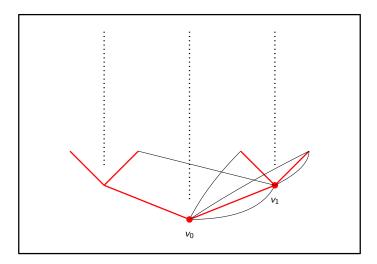
For  $(\mathbb{Q}, \leq)$  we have the Sierpiński colourings. Can we do something similar for the Rado graph?

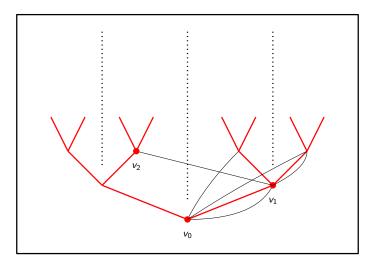


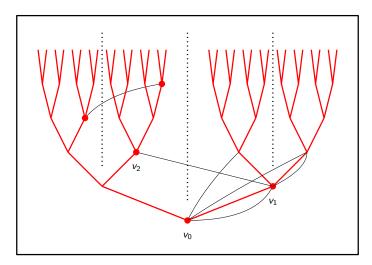


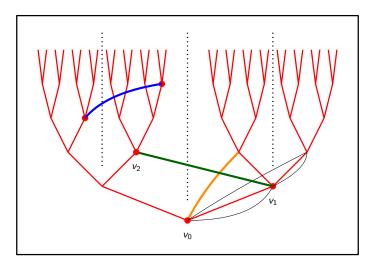


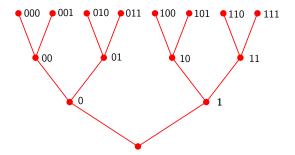






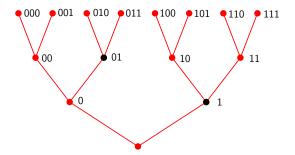






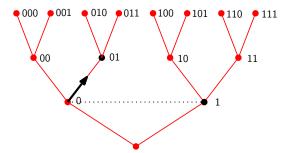
### Definition (Graph G)

- 1 Vertices:  $2^{<\omega}$
- 2 Vertices  $a, b \in 2^{<\omega}$  satisfying |a| < |b| forms and edge if and only if b(|a|) = 1.
- **3** There are no other edges.



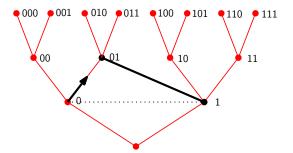
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## The upper bound

#### Lemma

**G** is universal: the Rado graph **R** embeds to **G**.

## Proof.

Assume that the vertex set of **R** is  $\omega$ . The vertex  $i \in \omega$  then corresponds to a sequence *a* of length *i* with a(j) = 1 if and only if  $i \sim j$ .

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The definition of **G** is stable for passing into a strong subtrees: if *S* is a strong subtree of  $2^{<\omega}$  then it is also a copy of **G** in **G** 

We thus can repeat precisely the same proof as before to obtain the upper bound on big Ramsey degrees.

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Lower bounds needs a bit more care.

## Thank you for the attention

Most we covered today is in S. Todorčević, Introduction to Ramsey spaces

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- C. Laflamme, L. Nguyen Van Thé, N. W. Sauer, Partition properties of the dense local order and a colored version of Milliken's theorem, Combinatorica 30(1) (2010), 83–104.